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# Chapter I

## Basic stuff

### 1 Categories

*September 16, 2023*

**Definition 1.1** (Category). A category  $\mathbf{C}$  consists of a class of objects  $\text{Obj } \mathbf{C}$ , and for every pair of objects  $X, Y$  of  $\mathbf{C}$ , a set of morphisms  $\text{Hom}_{\mathbf{C}}(X, Y)$  such that the following hold:

(i) (*Composition of morphisms*). For any objects  $X, Y, Z$  of  $\mathbf{C}$ , there's a function:

$$\begin{aligned} \text{Hom}_{\mathbf{C}}(X, Y) \times \text{Hom}_{\mathbf{C}}(Y, Z) &\rightarrow \text{Hom}_{\mathbf{C}}(X, Z) \\ (f, g) &\mapsto gf \end{aligned}$$

(ii) (*Associativity of composition*). For  $f \in \text{Hom}_{\mathbf{C}}(X, Y)$ ,  $g \in \text{Hom}_{\mathbf{C}}(Y, Z)$ ,  $h \in \text{Hom}_{\mathbf{C}}(Z, W)$ , we have

$$h(gf) = (hg)f.$$

(iii) (*Identity morphisms*). For any object  $X$  of  $\mathbf{C}$ , we have  $1_X \in \text{Hom}_{\mathbf{C}}(X, X)$ <sup>1</sup> such that for any  $f \in \text{Hom}_{\mathbf{C}}(X, Y)$  and for any  $g \in \text{Hom}_{\mathbf{C}}(Y, X)$ , we have

$$\begin{aligned} f 1_X &= f, \text{ and} \\ 1_X g &= g. \end{aligned}$$

(iv) (*Domains and codomains of morphisms*). For any objects  $X, Y, Z, W$  of  $\mathbf{C}$ , the sets  $\text{Hom}_{\mathbf{C}}(X, Y)$  and  $\text{Hom}_{\mathbf{C}}(Z, W)$  are disjoint unless  $X = Z$  and  $Y = W$ .

If  $\text{Obj}(\mathbf{C})$  forms a set, then  $\mathbf{C}$  is said to be *small*.

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<sup>1</sup>It's easily seen that such a morphism is unique for each  $X$ , justifying the notation  $1_X$ .

**Notation.** For a category  $\mathcal{C}$ , we'll also write " $f: X \rightarrow Y$  in  $\mathcal{C}$ " to mean that  $X, Y$  are objects in  $\mathcal{C}$  and  $f \in \text{Hom}_{\mathcal{C}}(X, Y)$ .

**Remark.** What we have defined are actually locally small categories. In general the hom-sets needn't be small. However, generally, a small category is one whose objects as well as morphisms form a set.

**Definition 1.2** (Subcategories). A category  $\mathcal{D}$  is said to be a subcategory of a category  $\mathcal{C}$  iff the following hold:

- (i)  $\text{Obj}(\mathcal{D})$  is a subclass of  $\text{Obj}(\mathcal{C})$ .
- (ii)  $\text{Hom}_{\mathcal{D}}(A, B) \subseteq \text{Hom}_{\mathcal{C}}(A, B)$  for all objects  $A, B$  of  $\mathcal{C}$ .
- (iii) The morphism composition in  $\mathcal{D}$  is inherited from that in  $\mathcal{C}$ .

If the " $\subseteq$ " in (ii) is "=", then  $\mathcal{D}$  is called a *full* subcategory of  $\mathcal{C}$ .

**Example 1.3** (Some examples).

- (i) (*Subcategories of Set*).  $\text{Set}$ ,  $\text{Top}$ ,  $\text{Grp}$ ,  $\text{Rng}$ ,  $\text{Ring}$ ,  $\text{Vect}_K$ ,  $\text{Mod}_R$ , their pointed versions.
- (ii) (*A category whose objects needn't be sets*). Let  $R$  be a relation on a set  $X$  which is reflexive and transitive. Then the elements of  $X$  form a category with the set of morphisms from  $a$  to  $b$  being

$$\begin{cases} \{(a, b)\}, & a R b \\ \emptyset, & \text{otherwise} \end{cases}$$

and the morphisms being given by

$$(b, c)(a, b) = (a, c).$$

- (iii) (*A category whose objects are sets but morphisms are not set theoretic functions*).  $\text{RelHTop}$ .
- (iv) (*Opposite category*). Let  $\mathcal{C}$  be a category. Then the objects of  $\mathcal{C}$  form another category  $\mathcal{C}^{\text{op}}$  with  $\text{Hom}_{\mathcal{C}^{\text{op}}}(X, Y) = \text{Hom}_{\mathcal{C}}(Y, X)$  and composition of  $f: X \rightarrow Y$  and  $g: Y \rightarrow Z$  in  $\mathcal{C}^{\text{op}}$  given by the composition of  $g$  and  $f$  in  $\mathcal{C}$ .

## 2 Monics, epics and isomorphisms

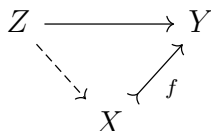
September 17, 2023

**Definition 2.1** (Monomorphisms and epimorphisms). Let  $f: X \rightarrow Y$  be a morphism in a category. Then:

- (i)  $f$  is called a monomorphism (or monic) iff

$$fu = fv \implies u = v$$

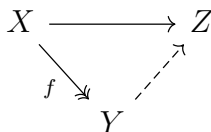
for any morphisms  $u, v: Z \rightarrow X$ , or equivalently, any morphism  $Z \rightarrow Y$  factors through  $f$  via at most one morphism:



- (ii)  $f$  is called an epimorphism (or epic) iff

$$uf = vf \implies u = v$$

for any morphisms  $u, v: Y \rightarrow Z$ , or equivalently, any morphism  $X \rightarrow Z$  factors through  $f$  via at most one morphism:



**Notation.** We'll sometimes denote monics by  $\succrightarrow$  and epics by  $\twoheadrightarrow$ .

**Definition 2.2** (Inverses and isomorphisms). A morphism  $f: X \rightarrow Y$  is called an isomorphism iff there exists another morphism  $g: Y \rightarrow X$  such that

$$\begin{aligned} gf &= 1_X, \text{ and} \\ fg &= 1_Y. \end{aligned}$$

Such a morphism is called an inverse of  $f$ , and is denoted by  $f^{-1}$ .<sup>2</sup>

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<sup>2</sup>Uniqueness of inverses (an easy fact) justifies this notation.

**Corollary 2.3.**

- (i) *Inverse of an isomorphism is an isomorphism.*
- (ii) *An isomorphism is monic and epic.*
- (iii) *In subcategory  $\mathbf{C}$  of  $\mathbf{Set}$ , injective morphisms are monic and surjective morphisms are epic. Further, if  $\mathbf{C}$  contains all the identity functions of its objects, then isomorphisms are bijective.*
- (iv) *Composition of monics (respectively epics) is monic (respectively epic).*
- (v) (a)  *$gf$  is monic  $\implies f$  is monic.*  
 (b)  *$gf$  is epic  $\implies g$  is epic.*

**Example 2.4** (Counters to the converse of (iii)).

- (i) (*Monic  $\not\Rightarrow$  injective*). In the category of “root-able” groups<sup>3</sup>,  $\mathbb{Q} \rightarrow \mathbb{Q}/\mathbb{Z}$  is monic.
- (ii) (*Epic  $\not\Rightarrow$  surjective*).
  - (a) In the category of Hausdorff spaces, any inclusion  $E \hookrightarrow X$  with  $E$  dense in  $X$ , is epic.<sup>4</sup>
  - (b) In the category of commutative rings with identities with homomorphisms preserving identities, the inclusion  $\mathbb{Z} \hookrightarrow \mathbb{Q}$  is epic.<sup>5</sup>

**Proposition 2.5.** *Monics are injective and epics, surjective in Grp.*

### 3 Initials, terminals, and zeroes

*September 24, 2023*

**Definition 3.1** (Initial, terminals, and zeroes). An object  $X$  in a category  $\mathbf{C}$  is called

- (i) *initial* iff  $\text{Hom}_{\mathbf{C}}(X, Y)$  is a singleton for each object  $Y$ ;
- (ii) *terminal* iff  $\text{Hom}_{\mathbf{C}}(Y, X)$  is a singleton for each object  $Y$ ; and,
- (iii) *zero* iff it is initial as well as terminal.

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<sup>3</sup>A group  $G$  is “root-able” iff  $G = G^n$  for every  $n \in \mathbb{Z} \setminus \{0\}$ . If  $G$  is abelian as well, we call it divisible, and in the additive notation, the condition reads  $G = nG$ . Fact: Any finite “root-able” group is trivial.

<sup>4</sup>More general: A continuous function on a Hausdorff codomain is determined by its restriction on a dense subset of its domain.

<sup>5</sup>More strongly, any homomorphism  $I \rightarrow R$  from an integral domain  $I$  extends uniquely to  $\text{Frac}(I) \rightarrow R$ .

**Corollary 3.2.**

- (i) *Initials (respectively terminals, zeroes) are unique up to unique isomorphisms.*
- (ii) *Any morphism into an initial object is epic and any morphism from a terminal object is monic.*

## 4 Zero morphisms

*September 24, 2023*

**Definition 4.1** (Left and right zeroes, and zero morphisms). A morphism  $f$  in a category is called a

- (i) *left zero* iff whenever defined,  $fu = fv$  for any morphisms  $u, v$ ;
- (ii) *right zero* iff whenever defined,  $uf = vf$  for any morphisms  $u, v$ ; and,
- (iii) *zero morphism* iff it's both, a left as well as a right zero.

**Remark.** “ $a$  followed by  $b$ ” will mean  $ba$  and not  $ab$ .

**Corollary 4.2.**

- (i) *The morphisms associated with initial (terminal) objects are right (left) zeroes.*
- (ii) *If  $f$  is a left zero, then  $fu$ , whenever defined, is too. Similarly for right zeroes.*
- (iii) *A right zero followed by a left zero is a zero morphism.*
- (iv) *If  $z$  is a zero morphism, then  $vzu$ , whenever defined is a zero morphism.*
- (v) *If  $0$  is a zero object, then  $X \rightarrow 0 \rightarrow Y$  is a zero morphism.*
- (vi) *In a category, there exists at most one family of zero morphisms between each pair of objects.*

**Definition 4.3** (A category having compatible zero morphisms). A category is said to have compatible zero morphisms iff there exists a family of zero morphisms  $0_{X,Y}: X \rightarrow Y$  for each pair of objects such that for any morphisms  $X \rightarrow Y$  and  $Y \rightarrow Z$ , the following diagram commutes:

$$\begin{array}{ccc}
 X & \xrightarrow{0_{X,Y}} & Y \\
 \downarrow & \searrow^{0_{X,Z}} & \downarrow \\
 Y & \xrightarrow{0_{Y,Z}} & Z
 \end{array}$$

**Corollary 4.4.** *If a category has a zero object  $0$ , then the morphisms  $X \rightarrow 0 \rightarrow Y$  form a compatible family of zero morphisms.*

## 5 Subobjects and quotients

September 24, 2023

**Definition 5.1** (Subobjects, quotients and their comparisons). In a category, *subobjects* of an object  $X$  are monics with codomain  $X$ , and *quotients* of  $X$  are epics with domain  $X$ .

If  $i_1: Y_1 \rightarrow X$  and  $i_2: Y_2 \rightarrow X$  are subobjects of  $X$ , then we write  $i_1 \leq i_2$  iff  $i_2$  factors through  $i_1$ :

$$\begin{array}{ccc} & X & \\ i_1 \nearrow & & \nwarrow i_2 \\ Y_1 & \dashrightarrow & Y_2 \end{array}$$

Similarly, for quotients  $q_1: X \rightarrow Z_1$  and  $q_2: X \rightarrow Z_2$  of  $X$ , we write  $q_1 \leq q_2$  iff  $q_2$  factors through :

$$\begin{array}{ccc} & X & \\ q_1 \searrow & & \swarrow q_2 \\ Z_1 & \dashleftarrow & Z_2 \end{array}$$

**Corollary 5.2.**

- (i) *If  $f: Y \rightarrow X$  is a subobject of  $X$  and  $g: Z \rightarrow Y$  a subobject of  $Y$ , then  $gf$  is a subobject of  $X$ . Similarly for quotients.*
- (ii) *The comparisons defined in Definition 5.1 form partial order with the “equality” replaced with “being isomorphic”.*

**Definition 5.3** (Images). An image of a morphism  $f: X \rightarrow Y$  is a smallest subobject of  $Y$  through which  $f$  factors:

$$\begin{array}{ccc} X & \xrightarrow{f} & Y \\ & \searrow & \nearrow \\ & \text{im } f & \\ & \vdots & \\ & Y' & \end{array}$$

(Note: Curved arrows from  $X$  and  $Y$  to  $Y'$  indicate that  $f$  factors through  $Y'$ .)

**Notation.** The arrow  $\text{im } f \rightarrow Y$  stands for any image of  $f$ , which are all isomorphic, due to Corollary 5.4.

**Corollary 5.4.** Images are unique up to unique isomorphisms.

## 6 (Co)equalizers and (co)products

October 25, 2023

**Remark.** Whenever the domain and codomain are clear from the context, we'll omit the subscript from  $0_{X,Y}$ .

**Definition 6.1** ((Co)equalizers and (co)kernels). An equalizer of  $f, g: X \rightarrow Y$  in a category is a morphism which is terminal among all the morphisms  $u: A \rightarrow X$  such that  $fu = gu$ :

$$\begin{array}{ccc} \text{eq}(f, g) & \longrightarrow & X \begin{array}{c} \xrightarrow{f} \\ \xrightarrow{g} \end{array} Y \\ & \swarrow \text{dashed} & \uparrow u \\ & & A \end{array}$$

If the category has zero morphisms, then an equalizer of  $f$  and  $0_{X,Y}$  is called a *kernel of  $f$* .

Coequalizers and *cokernels* are defined dually.

**Notation.** The arrow  $\text{eq}(f, g) \rightarrow X$  stands for any equalizer of  $f, g$ , which are all isomorphic due to Corollary 6.2 (ii). Similarly, we use  $\ker(f) \rightarrow X$ ,  $Y \rightarrow \text{coeq}(f, g)$ , and  $Y \rightarrow \text{coker}(f)$ .

**Corollary 6.2.**

- (i) Equalizers are monic and coequalizers, epic.
- (ii) (Co)equalizers are unique up to unique isomorphisms.
- (iii) (Characterizing (co)kernels). In a category with zero morphisms, a kernel  $\ker f \rightarrow X$  of  $f: X \rightarrow Y$  is characterized by being the terminal among all the morphisms  $u: A \rightarrow X$  such that  $fu = 0$ :

$$\begin{array}{ccc} & & 0 \\ & \text{curved arrow} & \\ \ker f & \longrightarrow & X \begin{array}{c} \xrightarrow{f} \\ \xrightarrow{0} \end{array} Y \\ & \swarrow \text{dashed} & \uparrow u \\ & & A \end{array}$$



*Dual characterization holds for cokernels.*

(iv) *Taking kernels of quotients, or cokernels of subobjects, reverses order.*

**Proposition 6.3** (Quotient topology as a coequalizer in  $\mathbf{Top}$ ). *Let  $X$  be a topological space and  $\sim$  an equivalence relation on  $X$  and  $R \subseteq X \times X$  be the corresponding subset. Consider  $X/\sim$  under quotient topology and  $R$  under the subspace topology. Then  $p: X \rightarrow X/\sim$  is a coequalizer of the projections  $R \rightarrow X$ .*

## 7 (Co)products

*December 10, 2023*

**Definition 7.1** ((Co)products). A product of objects  $X_i$ 's in a category is an object  $P$  together with morphisms  $\pi_i: P \rightarrow X_i$ 's which is terminal among all families of morphisms  $f_i: Y \rightarrow X_i$ :

$$\begin{array}{ccc} & X_i & \\ f_i \nearrow & & \nwarrow \pi_i \\ Y & \dashrightarrow & P \end{array}$$

Coproducts are defined dually.

**Corollary 7.2.** *(Co)products are unique up to unique isomorphisms.*

## 8 Abelian categories

*October 25, 2023*

**Definition 8.1** ((Pre-)additive and (pre-)abelian categories). A category  $\mathbf{C}$  is called:

- (i) *pre-additive* iff the following hold:
  - (a)  $\mathbf{C}$  has a zero object  $0$ .
  - (b) Each  $\mathrm{Hom}_{\mathbf{C}}(X, Y)$  forms an additive abelian group with  $0$  being the additive identity.
  - (c) Composition of morphisms is bilinear.
- (ii) *additive* iff it is pre-additive and has finite products and coproducts.
- (iii) *pre-abelian* iff it is additive and has kernels and cokernels.
- (iv) *abelian* iff it is pre-abelian, and all monics arise as kernels and epics as cokernels.

**Proposition 8.2.** *In a pre-additive category, the following hold:*

- (i) Finite products and coproducts coincide.<sup>6</sup>
- (ii)  $1_X$  is a kernel, and  $1_Y$  a cokernel of  $0_{X,Y}$ .
- (iii) For  $f: X \rightarrow Y$ ,
  - (a)  $0 \rightarrow X$  is a kernel of  $f \iff f$  is monic.
  - (b)  $Y \rightarrow 0$  is a cokernel of  $f \iff f$  is epic.
- (iv) A kernel of a cokernel of a kernel is a kernel, and a cokernel of a kernel of a cokernel is a cokernel.

**Corollary 8.3.** *In an abelian category, a monic is a kernel of any cokernel of itself, while an epic is a cokernel of any kernel of itself.*

**Proposition 8.4.** *In an abelian category, a morphism is an isomorphism  $\iff$  it is monic and epic.*

**Definition 8.5** (Exact sequences). In a pre-additive category, a sequence (finite or infinite at either end) of morphisms

$$\cdots \xrightarrow{f_{i-1}} X_{i-1} \xrightarrow{f_i} X_i \xrightarrow{f_{i+1}} X_{i+1} \longrightarrow \cdots$$

is said to be exact at a nonterminal object  $X_i$  iff kernels of  $f_{i+1}$  are precisely the images of  $f_i$ .

The sequence is called exact iff it is exact at all its nonterminal objects.

**Proposition 8.6** (Characterizing exactness). *In an abelian category, a sequence of morphisms*

$$A \xrightarrow{f} B \xrightarrow{g} C$$

*is exact  $\iff$  the following compositions are 0:*

$$\begin{array}{c} A \xrightarrow{f} B \xrightarrow{g} C \\ \ker g \longrightarrow B \longrightarrow \operatorname{coker} f \end{array}$$

**Corollary 8.7** (Using exactness to characterize various things). *In an abelian category, the exactness of the sequence of the left is equivalent to the property stated at*

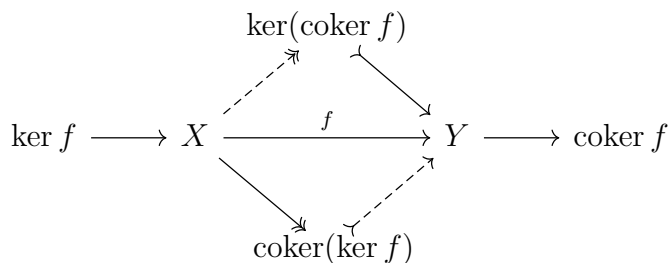
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<sup>6</sup>Only at the object level.

the right:

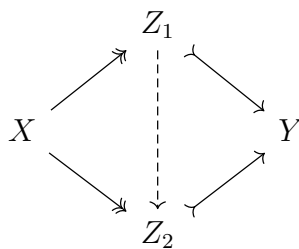
$$\begin{array}{ll}
 0 \longrightarrow X \xrightarrow{f} Y & f \text{ is monic} \\
 X \xrightarrow{f} Y \longrightarrow 0 & f \text{ is epic} \\
 0 \longrightarrow X \xrightarrow{f} Y \longrightarrow 0 & f \text{ is an isomorphism} \\
 0 \longrightarrow X \xrightarrow{f} Y \xrightarrow{g} Z & f \text{ is a kernel of } g \\
 X \xrightarrow{f} Y \xrightarrow{g} Z \longrightarrow 0 & g \text{ is a cokernel of } f \\
 0 \longrightarrow X \xrightarrow{f} Y \xrightarrow{g} Z \longrightarrow 0 & f \text{ is a kernel of } g \text{ and } g \text{ a cokernel of } f
 \end{array}$$

**Theorem 8.8.** *In an abelian category, any morphism  $f: X \rightarrow Y$  admits unique<sup>7</sup> factorizations as follows:*



**Proposition 8.9** ( $\text{im } f = \ker(\text{coker } f)$ ). *In an abelian category, a kernel of any cokernel of  $f$  is an image of  $f$ .*

**Corollary 8.10** (Unique factorization). *In an abelian category, the commutativity of solid arrows below implies the existence of the unique dashed arrow which makes the resulting diagram commute:*



Further, the dashed arrow is an isomorphism.

<sup>7</sup>Once  $\ker f$ ,  $\text{coker } f$ ,  $\text{coker}(\ker f)$  and  $\ker(\text{coker } f)$  have been fixed.

**Proposition 8.11** (Characterizing abelian-ness). *Let  $\mathcal{C}$  be a pre-abelian category. Then the following hold:*

(i) *For every  $f: X \rightarrow Y$ , there exists a unique<sup>8</sup>  $\tilde{f}$  such that the following diagram commutes.<sup>9</sup>*

$$\begin{array}{ccccc} \ker f & \longrightarrow & X & \longrightarrow & \operatorname{coker}(\ker f) \\ & & \downarrow f & & \downarrow \tilde{f} \\ \operatorname{coker} f & \longleftarrow & Y & \longleftarrow & \ker(\operatorname{coker} f) \end{array}$$

(ii)  $\mathcal{C}$  is abelian  $\iff \tilde{f}$  is an isomorphism for every  $f$ .

## 9 Functors

*December 10, 2023*

**Definition 9.1** (Functors). A functor  $F: \mathcal{C} \rightarrow \mathcal{D}$  is an assignment of objects and morphisms of  $\mathcal{C}$  to those in  $\mathcal{D}$  such that the following hold:

(i)  $F$  is compatible with domains and codomains:

$$\begin{array}{ccc} A & & F(A) \\ f \downarrow & \xrightarrow{F} & \downarrow F(f) \\ B & & F(B) \end{array}$$

(ii)  $F$  is compatible with composition, *i.e.*, commutativity of the diagram on the left implies that on the right:

$$\begin{array}{ccc} A & \xrightarrow{h} & C \\ f \searrow & & \nearrow g \\ & B & \end{array} \xrightarrow{F} \begin{array}{ccc} F(A) & \xrightarrow{F(h)} & F(C) \\ F(f) \searrow & & \nearrow F(g) \\ & F(B) & \end{array}$$

(iii)  $F$  is compatible with identities:

$$\begin{array}{ccc} A & & F(A) \\ 1_A \downarrow & \xrightarrow{F} & \downarrow 1_{F(A)} \\ A & & F(A) \end{array}$$

<sup>8</sup>Once  $\ker f$ ,  $\operatorname{coker} f$ ,  $\operatorname{coker}(\ker f)$  and  $\ker(\operatorname{coker} f)$  have been fixed.

<sup>9</sup> $\operatorname{coker}(\ker f)$  is shortened for  $\operatorname{coker}(\ker f \rightarrow X)$ . Similarly for  $\ker(\operatorname{coker} f)$ .

**Example 9.2** (Hom-functors). Let  $A, B$  be objects of a category  $\mathcal{C}$ . Then the following defines the functor  $\text{Hom}_{\mathcal{C}}(A, -): \mathcal{C} \rightarrow \text{Set}$ :

$$\begin{array}{ccc} X & & \text{Hom}_{\mathcal{C}}(A, X) \\ f \downarrow & \longmapsto & \downarrow_{\text{Hom}_{\mathcal{C}}(A, f)} \\ Y & & \text{Hom}_{\mathcal{C}}(A, Y) \end{array}$$

where  $\text{Hom}_{\mathcal{C}}(A, f)$  is given by  $\alpha \mapsto f\alpha$ .

We also have the functor  $\text{Hom}_{\mathcal{C}}(-, B): \mathcal{C}^{\text{op}} \rightarrow \text{Set}$  given by:<sup>10</sup>

$$\begin{array}{ccc} X & & \text{Hom}_{\mathcal{C}}(X, B) \\ f \downarrow & \longmapsto & \uparrow_{\text{Hom}_{\mathcal{C}}(f, B)} \\ Y & & \text{Hom}_{\mathcal{C}}(Y, B) \end{array}$$

where  $\text{Hom}_{\mathcal{C}}(g, B)$  is given by  $\beta \mapsto \beta f$ .<sup>11</sup>

**Corollary 9.3.**

(i) *Functors preserve isomorphisms.*

(ii) (Composition of functors). *Let  $F: \mathcal{C} \rightarrow \mathcal{D}$  and  $G: \mathcal{D} \rightarrow \mathcal{E}$  be functors. Then the following defines their composite functor  $GF: \mathcal{C} \rightarrow \mathcal{E}$ :*

$$\begin{array}{ccc} A & & G(F(A)) \\ f \downarrow & \xrightarrow{GF} & \downarrow_{G(F(f))} \\ B & & G(F(B)) \end{array}$$

(iii) (Identity functors). *For any category  $\mathcal{C}$ , we have an identity functor  $1_{\mathcal{C}}$ :*

$$\begin{array}{ccc} A & & A \\ f \downarrow & \xrightarrow{1_{\mathcal{C}}} & \downarrow_f \\ B & & B \end{array}$$

(iv) (Category of categories). *Small categories with functors as morphisms form a category  $\text{Cat}$ .*

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<sup>10</sup>Note that  $f: X \rightarrow Y$  is in  $\mathcal{C}$ .

<sup>11</sup>The composition is taking in  $\mathcal{C}$ .

**Definition 9.4** (Isomorphism of categories). A functor  $F: \mathbf{C} \rightarrow \mathbf{D}$  is called an isomorphism iff there exists another functor  $G: \mathbf{D} \rightarrow \mathbf{C}$  such that  $GF = 1_{\mathbf{C}}$  and  $FG = 1_{\mathbf{D}}$ .<sup>12</sup>

**Definition 9.5** (Full, faithful and essentially surjective functors). A functor  $F: \mathbf{C} \rightarrow \mathbf{D}$  is called *full* (respectively *faithful*) iff for all objects  $A, B$  of  $\mathbf{C}$ , the associated function  $\text{Hom}_{\mathbf{C}}(A, B) \rightarrow \text{Hom}_{\mathbf{D}}(F(A), F(B))$  is surjective (respectively injective).

$F$  is called *essentially surjective* iff for every object  $X$  of  $\mathbf{D}$ , there exists an object  $A$  of  $\mathbf{C}$  such that  $F(A)$  is isomorphic to  $X$ .

**Lemma 9.6** (Fully faithful functors are conservative). *Let  $F: \mathbf{C} \rightarrow \mathbf{D}$  be fully faithful and  $f$  be a morphism in  $\mathbf{C}$ . Then  $f$  is an isomorphism  $\iff F(f)$  is an isomorphism.*

## 10 Natural transformations

*December 11, 2023*

**Definition 10.1** (Natural transformations). Let  $F, G: \mathbf{C} \rightarrow \mathbf{D}$  be functors. Then a natural transformation  $\eta: F \Rightarrow G$  is a family of morphisms  $\eta_A$  in  $\mathbf{D}$  indexed by the objects  $A$  on  $\mathbf{C}$  such that the following hold:

- (i)  $\eta_A: F(A) \rightarrow G(A)$  in  $\mathbf{D}$ .
- (ii) For any morphism  $f: A \rightarrow B$  in  $\mathbf{C}$ , the following diagram commutes:

$$\begin{array}{ccc} F(A) & \xrightarrow{\eta_A} & G(A) \\ F(f) \downarrow & & \downarrow G(f) \\ F(B) & \xrightarrow{\eta_B} & G(B) \end{array}$$

If each  $\eta_A$  is an isomorphism, then  $\eta$  is called a natural *isomorphism*.

**Notation.** We'll denote  $\text{Nat}(F, G)$  to denote the class of all natural transformations  $F \Rightarrow G$ .

**Corollary 10.2.**

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<sup>12</sup>Clearly, for small categories, isomorphism of categories is precisely an isomorphism between them in  $\text{Cat}$ .

(i) (Composition of natural transformations). Let  $F, G, H: \mathbf{C} \rightarrow \mathbf{D}$  be functors and  $\eta: F \Rightarrow G$  and  $\xi: G \Rightarrow H$  be natural transformations. Then the following defines the composite natural transformation  $\xi\eta: F \Rightarrow H$ :

$$(\xi\eta)_A := \xi_A \eta_A$$

(ii) (Identity natural transformations). Let  $F: \mathbf{C} \rightarrow \mathbf{D}$  be a functor. Then the following defines the identity natural transformation  $1_F: F \Rightarrow F$ :

$$(1_F)_A := 1_{F(A)}$$

(iii) (Functor category). Let  $\mathbf{C}$  be a small category. Then for any category  $\mathbf{D}$ , the functors  $\mathbf{C} \rightarrow \mathbf{D}$  form a category  $\text{Func}(\mathbf{C}, \mathbf{D})$  with the morphisms being the natural transformations between them.

Further, a natural isomorphism between functors  $\mathbf{C} \rightarrow \mathbf{D}$  is precisely an isomorphism in  $\text{Func}(\mathbf{C}, \mathbf{D})$ .

**Definition 10.3** (Equivalence of categories). A functor  $F: \mathbf{C} \rightarrow \mathbf{D}$  is called an equivalence iff there exists another functor  $G: \mathbf{D} \rightarrow \mathbf{C}$  such that  $GF$  is naturally isomorphic to  $1_{\mathbf{C}}$  and  $FG$  to  $1_{\mathbf{D}}$ .

**Theorem 10.4** (Characterizing equivalence). Let  $F: \mathbf{C} \rightarrow \mathbf{D}$  be a functor. Then  $F$  is an equivalence  $\iff F$  is fully faithful and essentially surjective.<sup>13</sup>

## 11 Yoneda lemma

December 11, 2023

**Proposition 11.1** (Yoneda lemma). Let  $\mathbf{C}$  be a small category.<sup>14</sup> Let  $F: \mathbf{C}^{\text{op}} \rightarrow \text{Set}$  be a functor and  $B \in \text{Obj}(\mathbf{C})$ . Then the function  $\text{Nat}(\text{Hom}_{\mathbf{C}}(-, B), F) \rightarrow F(B)$  given by

$$\eta \mapsto \eta_B(1_B)$$

is a bijection with the inverse given by

$$c \mapsto \left\{ \begin{array}{l} \text{Hom}_{\mathbf{C}}(X, B) \rightarrow F(X) \\ f \mapsto F(f)(c) \end{array} \right\}_{X \in \text{Obj}(\mathbf{C})}$$

<sup>13</sup>“ $\Leftarrow$ ” requires choice for classes.

<sup>14</sup>For  $\text{Nat}(\text{Hom}_{\mathbf{C}}(-, Y), F)$  to be a set.

**Lemma 11.2** (The canonical functor  $\mathbf{C} \rightarrow \mathbf{Funct}(\mathbf{C}^{\text{op}}, \mathbf{Set})$ ). *Let  $\mathbf{C}$  be a small category.<sup>15</sup> Then the following defines a functor  $\mathbf{C} \rightarrow \mathbf{Funct}(\mathbf{C}^{\text{op}}, \mathbf{Set})$ :*

$$\begin{array}{ccc} A & & \text{Hom}_{\mathbf{C}}(-, A) \\ f \downarrow & \longmapsto & \Downarrow \eta(f) \\ B & & \text{Hom}_{\mathbf{C}}(-, B) \end{array}$$

where  $\eta(f)$  is the natural transformation given by

$$\begin{aligned} \eta(f)_X: \text{Hom}_{\mathbf{C}}(X, A) &\longrightarrow \text{Hom}_{\mathbf{C}}(X, B) \\ \alpha &\longmapsto f\alpha. \end{aligned}$$

**Proposition 11.3** (Yoneda theorem). *Let  $\mathbf{C}$  be a small category. Then the canonical functor  $\mathbf{C} \rightarrow \mathbf{Funct}(\mathbf{C}^{\text{op}}, \mathbf{Set})$  is fully faithful.*

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<sup>15</sup>So that we can talk about  $\mathbf{Funct}(\mathbf{C}^{\text{op}}, \mathbf{Set})$ .